

Meta-Predicates for Rodin

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Rodin Workshop 2016

- Common Patterns in Event-B
 - PO for deadlock freedom
 - PO for deterministic behavior
- Realized by composing guards
 - in other guards
 - in invariants
- Can be automated!

Deadlock Freedom - Example

```
event evt1
any
  x1
where
  x1 > 4
then
  counter := counter + x1
end

event evt2
any
  x2
where
  x2 < 2
then
  counter := counter - x2
end
```

Deadlock Freedom $\Leftrightarrow \exists x1 . x1 > 4 \vee \exists x2 . x2 < 2$

The Problem

Cut	⌘X
Copy	⌘C
Paste	⌘V
Paste and Match Style	⇧⌘V

But we know how to do it!

We found copy and paste errors in the ABZ'14 case studies.

Meta-Predicate in Invariant

EXTENDED INVARIANTS

- $\neg(\text{enabled}(\{\text{error_case_event}\}))$ not theorem >

INVARIANTS

- $\text{inv1: } \text{pc} \in \mathbb{N}$ theorem >

More Involved Invariant

EXTENDED INVARIANTS

- `enabled({recursive_case}) ⇒ ¬enabled({recursive_case_fail})` not theorem >

Matching Guard

```
recursive_case_fail:      not extended ordinary >  
REFINES  
o   terminate_with_ce  
ANY  
o   violating_state >  
EXTENDED GUARDS  
o    $\neg$  enabled({recursive_case}) not theorem >  
WHERE  
o   grd1:   pc = recursive_case not theorem >  
o   find_violation: PROPERTY(violating_state) = FALSE not theorem >  
THEN  
o   act1:   result = counterexample_found >  
o   act2:   counterexample=violating_state >  
END
```


Result

- ▼ ⓘ Proof Obligations
 - ⓘ generated_invariant_2/WD
 - ⓘ INITIALISATION/generated_invariant_2/INV
 - ✓ basecase_init_fail/grd2/WD
 - ✓^A basecase_init_fail/find_violation/GRD
 - ✓ basecase_init/grd2/WD
 - ⓘ basecase_init/generated_invariant_2/INV
 - ✓ basecase_transition_fail/grd2/WD
 - ✓^A basecase_transition_fail/find_violation/GRD
 - ✓ basecase_transition/grd2/WD
 - ⓘ basecase_transition/generated_invariant_2/INV

Replaced by static checker, completely transparent to Rodin Platform

New Predicates

- *deadlock*, all events disabled

$$\text{deadlock}(\text{Events}) = \bigwedge_{e \in \text{Events}} \bigvee_{g \in \text{guards}(e)} \neg g$$

- *enabled*, all events enabled

$$\text{enabled}(\text{Events}) = \bigwedge_{e \in \text{Events}} \bigwedge_{g \in \text{guards}(e)} g$$

New Predicates Continued

- *controller*, exactly one enabled

$$\text{controller}(\text{Events}) = \bigvee_{e \in \text{Events}} (\text{enabled}(\{e\}) \wedge \text{deadlock}(\text{Events} \setminus \{e\}))$$

- *deterministic* order of execution

$$\text{deterministic}(\text{Events}) = \text{controller}(\text{Events}) \vee \text{deadlock}(\text{Events})$$

Problematic Example

```
event evt1
any
  x1
where
  x1 > 4
then
  counter := counter + x1
end
```

```
event evt2
any
  x2
where
  enabled(evt1)
  x2 < 2
then
  counter := counter - x2
end
```

Problematic Example

```
event evt1
any
  x1
where
  x1 > 4
then
  counter := counter + x1
end
```

```
event evt2
any
  x1, x2
where
  x1 > 4
  x2 < 2
then
  counter := counter - x2
end
```

```
event evt2
any
  x2
where
  #(x1).(x1 > 4)
  x2 < 2
then
  counter := counter - x2
end
```

Yet Another Problematic Example

```
event evt1
any
  x1
where
  x1 > 4
then
  counter := counter + x1
end

event evt2
any
  x1
where
  enabled(evt1)
  x1 < 2
then
  counter := counter - x1
end
```

Problematic Example

```
event evt1
any
  x1
where
  x1 > 4
then
  counter := counter + x1
end
```

```
event evt2
any
  x1
where
  x1 > 4
  x1 < 2
then
  counter := counter - x1
end
```

```
event evt2
any
  x1
where
  #(x1).(x1 > 4)
  x1 < 2
then
  counter := counter - x1
end
```

Decide on Scoping

```
event evt1
any
  x1
where
  x1 > 4
then
  counter := counter + x1
end
```

```
event evt2
any
  x1
where
   $\neg \text{enabled}(\text{evt1})$ 
  x1 > 2
then
  counter := counter - x1
end
```


Decide on Scoping

```
event evt1
any
  x1
where
  x1 > 4
then
  counter := counter + x1
end
```

```
event evt2
any
  x1
where
   $\neg x1 > 4$ 
  x1 > 2
then
  counter := counter - x1
end
```

```
event evt2
any
  x1
where
   $\neg \exists x1. x1 > 4$ 
  x1 > 2
then
  counter := counter - x1
end
```

Decide on Scoping

- Always introduce using quantifier
- Only introduce if new variable name
- Introduce using parameter
- Introduce preference

Conclusions

- Easy to implement using Rodin's API
- Less error-prone
 - No copy & paste
 - No missing updates
 - Expressiveness
- Useful for analysis:
 - Static checking
 - Constraint based verification
 - Model checking
 - Control flow
 - ...

- Sources and Bugtracker
`github.com/wysiib/RodinMetaPredicatesPlugin`
- Update Site for Rodin
`www3.hhu.de/stups/rodin/meta_predicates/nightly`
- Currently only nightly builds

Thank you!