



A (very) small experiment in Event-B rippling

Gudmund Grov, Alan Bundy & Lucas Dixon

Rodin Workshop, Dusseldorf 2010

Talk outline

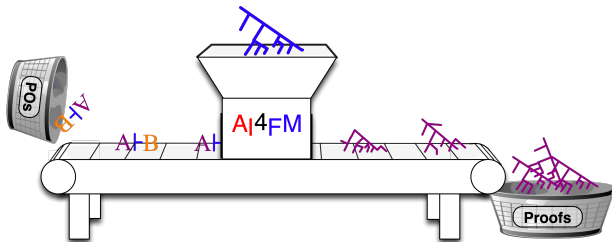
- ▶ “Abrial’s MMPE rule” : $n * x/100 * f * p * 20$
 - ▶ 100,000 loc \rightsquigarrow 3-12 Man Months of Proof Effort
- ▶ BUT, \exists families of Event-B UPOs based on proof strategy
 - ▶ AI4FM tries to explore such families to increase automation.
- ▶ This requires high-level proof strategies
 - ▶ rippling is an example of a high level proof strategy
 - ▶ implemented in Isaplanner.
- ▶ In this talk we will:
 - ▶ give an overview of the AI4FM project
 - ▶ describe a simple Event-B experiments with rippling/Isaplanner
 - ▶ beg for help!

AI4FM overview

- ▶ The user manually proves one exemplar proof.



- ▶ The theorem prover uses the additional information from the exemplar proof to discharge “similar” proofs:



The project

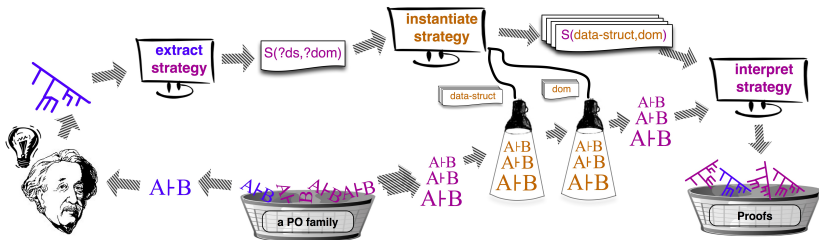
- ▶ 4 years UK EPSRC funding – started 1 April 2010
 - ▶ EP/H024050/1, EP/H024204/1 and EP/H023852/1.
- ▶ The team
 - ▶ **Newcastle**
 - ▶ Cliff Jones, Leo Freitas, Andrius Velykis
 - ▶ **Edinburgh**
 - ▶ Alan Bundy, Gudmund Grov, Yuhui Lin
 - ▶ **Heriot-Watt**
 - ▶ Andrew Ireland
 - ▶ **Southampton**
 - ▶ Michael Butler.

Nature of Event-B POs/proofs

- ▶ Rarely deep.
 - ▶ we are *not* trying to prove real math theorems!
- ▶ Complexity reduced by layering abstractions.
- ▶ Lots of POs
 - ▶ ... which can often be grouped into “(proof) families”.
- ▶ Lots of detail (on larger examples)

The AI4FM process

- ▶ (Somehow) classify POs into families.
- ▶ Require the expert user to prove 1 PO manually/interactively
 - ▶ preferable the “simplest”.
- ▶ This extra information is used to discharge rest of families
- ▶ Requires abstracting proof into a higher-level strategy
 - ▶ ... and we must design a strategy language to capture this.



Towards a strategy language

- ▶ The strategy language needs to be robust to cope with changes.
- ▶ A strategy may describe a sequence of intermediate lemmas that could be spawned:
 prove lemmas L_1, L_2, \dots, L_n s.t. GOAL follows.
- ▶ Abstract over many “dimensions”
 - ▶ data-structure, domain, etc.
- ▶ Include notions like generalisation and lemma discovery.
- ▶ Rippling provides evidence for a high-level strategy language.

Rippling (in a hurry)

- ▶ Proof plans: high-level description of proofs
 - ▶ captures common patterns of reasoning
- ▶ Rippling: a proof plan, which
 - ▶ works when one of the givens can be embedded in the goal (e.g inductive step cases)
 - ▶ for example in an Event-B INV type PO:

$$com = Cls \triangleleft (cl; nm) \vdash com = Cls \triangleleft ((cl \triangleleft \{s \mapsto (size \mapsto [])\})^\uparrow ; nm)$$

- ▶ annotations to guide rewriting (towards given)
 - ▶ language of **wave-fronts** and **skeletons**
- ▶ (direction) guarantees termination

Rippling illustration (step case of $t@(Y@Z) = (t@Y)@Z$)

Rewrite rules:

$$H\#T^{\uparrow} @L \Rightarrow H\#T@L^{\uparrow} \quad (1)$$

$$X_1\#X_2^{\uparrow} = Y_1\#Y_2^{\uparrow} \Rightarrow X_1 = Y_1 \wedge X_2 = Y_2^{\uparrow} \quad (2)$$

Rippling proof:

$$t @ (Y @ Z) = (t @ Y) @ Z \quad \text{IH}$$

$$h\#t^{\uparrow} @ (Y @ Z) = (h\#t^{\uparrow} @ Y) @ Z \quad \text{apply (1) 2 times}$$

$$h\#t@(Y@Z)^{\uparrow} = h\#(t@Y)^{\uparrow} @ Z \quad \text{apply (1)}$$

$$h\#t@(Y@Z)^{\uparrow} = h\#(t@Y)@Z^{\uparrow} \quad \text{apply (2)}$$

$$h = h \wedge t@(Y@Z) = (t@Y)@Z^{\uparrow} \quad \text{apply IH}$$

IsaPlanner

- ▶ Isaplanner is a proof planner built on top of Isabelle
 - ▶ enables use of existing Isabelle automation
 - ▶ soundness ensured by Isabelle.
- ▶ Reasoning techniques which generates proof plans
 - ▶ lazy search over possible ways to apply technique
- ▶ Implemented in the ML language
 - ▶ language also used to write new reasoning techniques
 - ▶ provides many ML functions to aid developing new techniques
- ▶ Rippling technique encoded (among others)

The experiment

- ▶ Goal is to check how rippling works in an Event-B setting
- ▶ Longer term we hope to see how easy it is to encode new techniques/patterns in Isaplanner.
- ▶ Set representation in Isabelle/HOL used
 - ▶ ... and slightly extended
- ▶ Modelled an example system in Rodin
 - ▶ translated POs into Isabelle/HOL representation
- ▶ Addressed INV type POs

Representing Event-B POs in Isabelle/HOL

- ▶ Event-B POs (+ theory) must be represented in Isabelle/HOL
- ▶ Uses Isabelle/HOL's set theory
- ▶ Extended with “Event-B operators”
 - ▶ e.g. \triangleleft , \triangleright , $\triangleleft\!\triangleright$, $\triangleright\!\triangleleft$, ...
- ▶ Functions as relations (HOL functions are total)
 - ▶ \rightarrow , \mapsto , \rightsquigarrow , ... defined
 - ▶ function application uses definite description operator (i.e. ι)
- ▶ Proof rules (theorems) derived by need
- ▶ Drawbacks
 - ▶ Most general type for functions
 - ▶ little mileage from HOL type checking
 - ▶ WDs ignored
 - ▶ Ignores Event-B proof rules

The system: a telephone system

- ▶ Based on Z model by Woodcock
- ▶ The variables
 - ▶ $call \in Subs \leftrightarrow (Status \leftrightarrow Subs)$
 - ▶ $connected \in Subs \leftrightarrow Subs$
 - ▶ st and num projections on $(Status \leftrightarrow Subs)$
 - ▶ $Free \subseteq Subs$

- ▶ Invariants

$inv1 : Callers = dom((call; st) \triangleright Connected)$

$inv2 : connected = Callers \triangleleft (call; num)$

- ▶ where $Connected \subset Status$ for "connected calls".

The system: a telephone system (events)

- ▶ Event's to lift handle, dial, answer, etc..
- ▶ We will only discuss two events

- ▶ **EVENT** *LiftFree* $\hat{=}$

- ANY** s **WHERE** $s \in Free$

- THEN** $Free := Free \setminus \{s\}$

- $call(s) := (seize \mapsto empty)$

- ...

- ▶ **EVENT** *LiftSuspended* $\hat{=}$

- ANY** s **WHERE** $(s \mapsto suspended) \in connected^{-1}; call; st$

- THEN** $call(connected^{-1}(s)) := (speech \mapsto s)$

- ...

PO LiftFree/inv1/INV

- ▶ All required rewrite (wave) rules are existing
- ▶ Proof strategy
 - ▶ Rippling followed by application of IH (inv1)
 - ▶ Then manually discharge reminding (non-rippling) goals
 - ▶ i.e. the conditions from conditional rewrite rules
- ▶ PO LiftFree/inv1/INV:

$$\begin{aligned} \text{Callers} &= \text{dom}(\text{call}; st \triangleright \text{Connected}), s \in \text{Free} \vdash \\ &\text{Callers} = \text{dom}((\text{call} \triangleleft \{s \mapsto (\text{seize} \mapsto \text{empty})\}); st \triangleright \text{Connected}) \end{aligned}$$

- ▶ (in one branch) gives the following two (provable) sub-goals

$$s \notin \text{dom}(\text{call}) \quad ((s \mapsto (\text{seize} \mapsto [])); st) \triangleright \text{Connected} = \{\}$$

- ▶ note there are also several (sometimes unprovable) branches

PO LiftFree/inv2/INV

- ▶ PO LiftFree/inv2/INV:

$$\begin{aligned} \text{connected} &= \text{Callers} \triangleleft (\text{call}; \text{num}), s \in \text{Free} \vdash \\ &\text{connected} = \text{Callers} \triangleleft ((\text{call} \triangleleft \{s \mapsto (\text{seize} \mapsto \text{empty})\}); \text{num}) \end{aligned}$$

- ▶ (in one branch) gives the following two (provable) sub-goals

$$s \notin \text{dom}(\text{call}) \quad (\text{Callers} \triangleleft \{(s \mapsto (\text{seize} \mapsto []))\}; \text{num}) = \{\}$$

- ▶ The “idea” is similar to LiftFree/inv1/INV
 - ▶ “ripple out” to “isolate” the new part – and show it is $\{\}$
 - ▶ BUT the rules used, i.e. the proof, differs (e.g. \triangleright vs. \triangleleft)
 - ▶ AND one of the sub-goals is harder to prove

LiftSuspended POs

▶ PO LiftSuspended/inv1/INV

$Callers = dom(call; st \triangleright Connected),$
 $(s \mapsto suspended) \in connected^{-1}; call; st \vdash$
 $Callers = dom((call \leftarrow \{connected^{-1}(s) \mapsto (speech \mapsto s)\}); st \triangleright Connected)$

▶ PO LiftSuspended/inv2/INV

$connected = Callers \triangleleft (call; num),$
 $(s \mapsto suspended) \in connected^{-1}; call; st \vdash$
 $connected = Callers \triangleleft ((call \leftarrow connected^{-1}(s) \mapsto (speech \mapsto s)); num)$

- ▶ Previous approach: s was not in $Callers$, and is still not
- ▶ Here, the “dual” is true:
 - ▶ $connected^{-1}(s)$ was in $Callers$ and still is.
- ▶ Requires “theory” properties of num , st , $call$ and $connected$

Some observations: difference with existing rippling work

- ▶ Hard to make any conclusions from a small case study
- ▶ Rippling seems promising for INV POs
- ▶ It manages to reduce the PO to a smaller, simpler goal – however:
 - ▶ much more conditional rewriting than other domains
 - ▶ .. and WD (which will add further conditions) ignored
 - ▶ not clear how to discharge these (non-rippling) goal
- ▶ Still lots of room for improvement:
 - ▶ several branches – use counter-example finders to filter out most obvious ones
 - ▶ better use of existing techniques (simp, blast, etc)
 - ▶ productive use of failure.

Productive use of failure

- ▶ One advantage of rippling is Ireland's proof critics.
- ▶ Particular failures (or partial matches) of rippling triggers certain "exception cases"
 - ▶ e.g. a missing lemma is speculated, or the conjecture is generalised.
- ▶ More robust than lower-level tactics.
- ▶ These could be applied in the Event-B setting
 - ▶ but not clear how to prove discovered/generalised lemmas
 - ▶ induction+rippling traditionally used
 - ▶ ... this may not be the case in Event-B
 - ▶ we may need to discover new proof strategies here
 - ▶ ... or learn them from a given exemplar proof ...

Rippling experiment vs. AI4FM agenda

- ▶ AI4FM is about learning strategies from an exemplar proof
- ▶ Rippling is an existing strategy
- ▶ Maybe a general strategy for INV is to use rippling
- ▶ and the special purpose strategies learned from exemplar proofs are used to
 - ▶ discharge sub-goals from conditional rewrite rules in target proofs
 - ▶ .. and (internally) of source proof, use proof of one such sub-goal to discharge other
 - ▶ prove discovered lemmas.

The AI4FM plan

- ▶ **Analyse** a lot of
 - ▶ POs
 - ▶ preferably from real-world applications
 - ▶ ... their (expert-provided) proofs attempts
 - ▶ and analysing “families” .
- ▶ **Based on analysis, *develop a strategy language***
 - ▶ more **examples** ⇒ more **robustness**
 - ▶ thus, the need for **examples!**
 - ▶ we expect an **iterative** development of the language.
- ▶ **Provide tool support**
 - ▶ to extract a strategy out of an exemplar proof
 - ▶ to interpret strategies to discharge “similar POs” .

We need your help (so we can help you)!

- ▶ We have experimented with our own little case-study.
- ▶ Steve Wright have already volunteered to participate!
- ▶ You can help by providing case-studies with (non-trivial) POs:



- ▶ ... industrial-sized preferably
 - ▶ ideally, with (given) families of proofs
 - ▶ even better with with proof history (including dead-ends)
 - ▶ the full proof process can tell us more than the finished article
-
- ▶ Examples from difference sources will increase robustness
 - ▶ ... which will (hopefully) give you **better proof automation!**

Summary

- ▶ AI4FM is a new project trying to learn and recycle proof strategies.
- ▶ Rippling is a high-level proof strategy
- ▶ We have shown promising results when rippling Event-B POs
 - ▶ a simple example
 - ▶ user-interaction still required
- ▶ (Your?) relevant example(s) will be of great help!

Thank you!



<http://www.ai4fm.org>