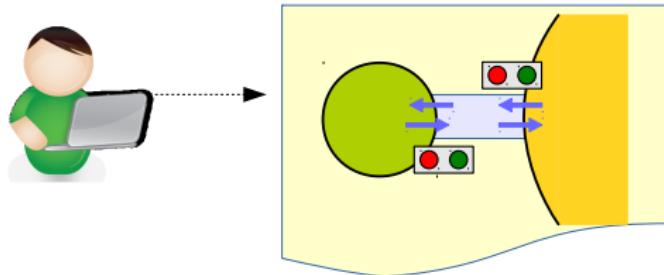


# SliceAndMerge: A Rodin Plug-in for Refactoring Refinement Structure of Event-B Machines

Tsutomu Kobayashi (University of Tokyo),  
Aivar Kripsaar (RWTH Aachen University),  
Fuyuki Ishikawa (NII, Japan),  
and Shinichi Honiden (NII, Japan)

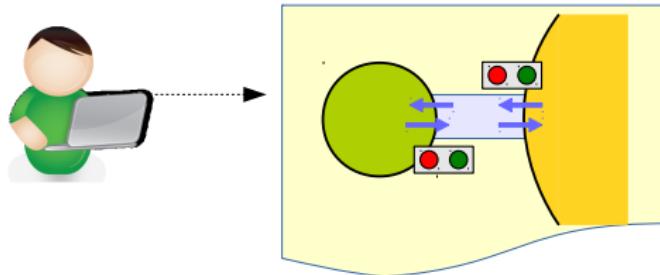
Rodin Workshop 2016  
May 23, 2016

# Modeling in Event-B

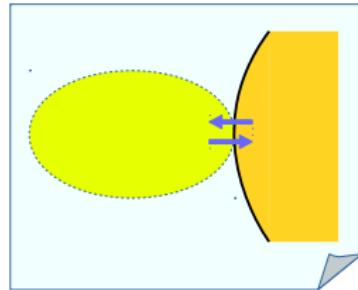


Informal document of  
traffic control system

# Modeling in Event-B

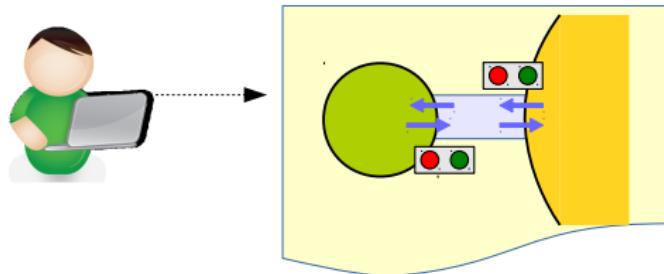


Informal document of traffic control system

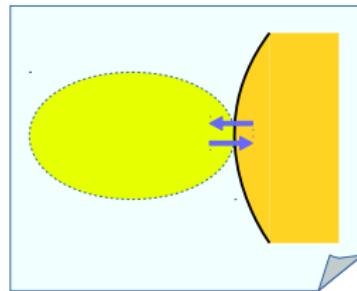


Abstract machine

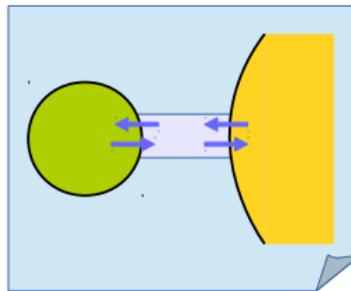
# Modeling in Event-B



Informal document of  
traffic control system

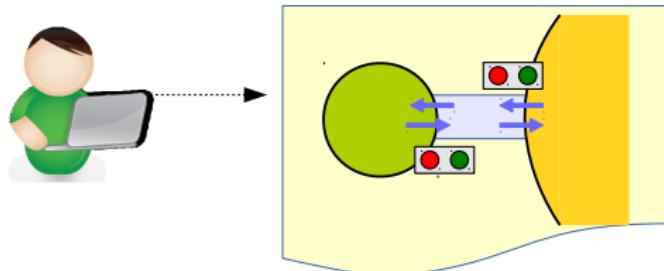


Abstract machine



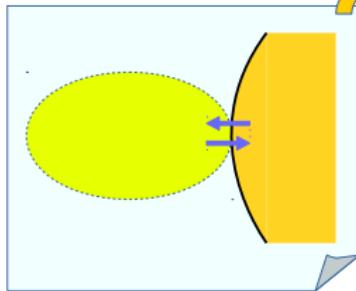
Medium machine

# Modeling in Event-B

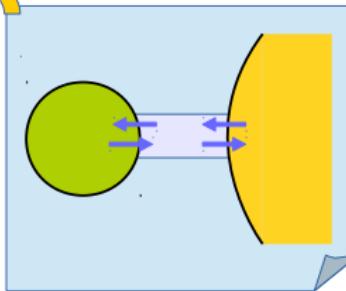


Informal document of  
traffic control system

prove  
Consistent?

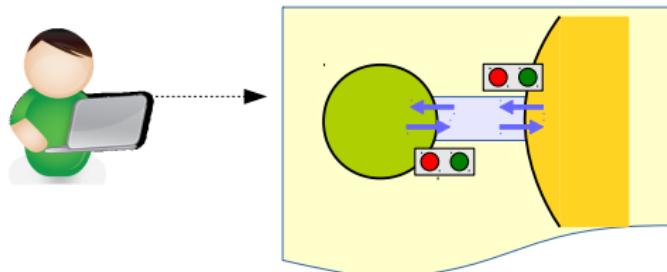


Abstract machine



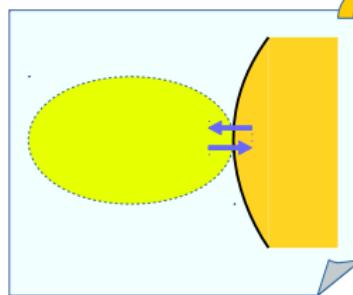
Medium machine

# Modeling in Event-B

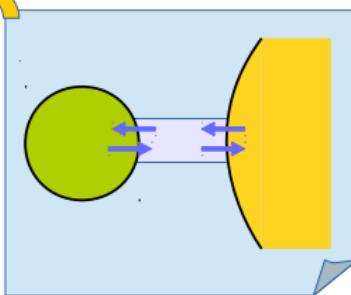


Informal document of  
traffic control system

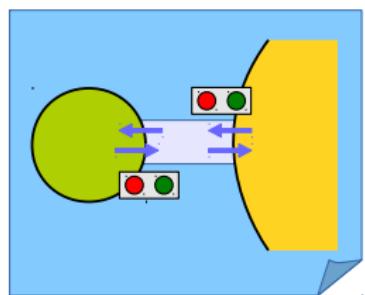
prove  
Consistent?



Abstract machine

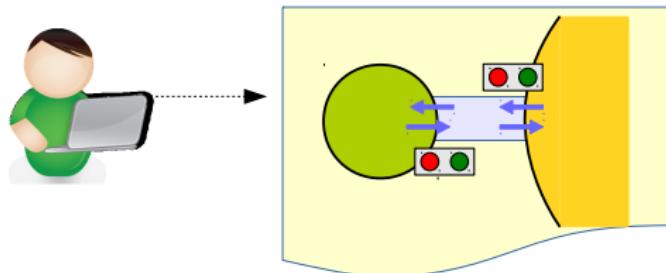


Medium machine



Concrete machine

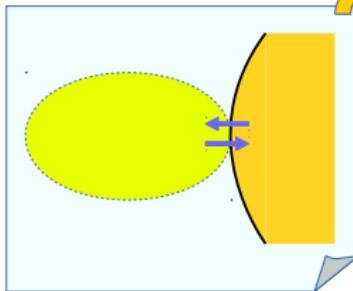
# Modeling in Event-B



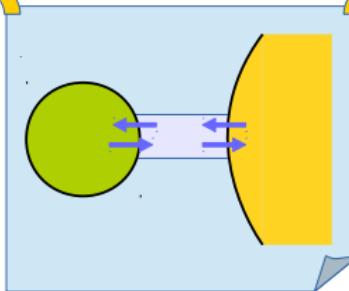
Informal document of  
traffic control system

prove  
Consistent?

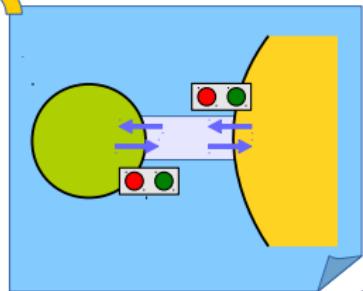
prove  
Consistent?



Abstract machine



Medium machine



Concrete machine

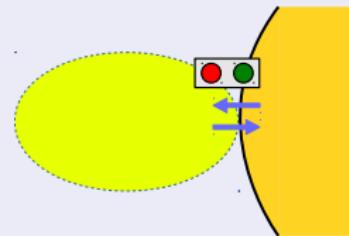
# Motivation

## Refinement design and specifications

Designing refinement = designing target system's aspects of interest

e.g., Focus on

- the outside
- traffic lights on the mainland



Specify/verify properties of them

- *If traffic light is green,  
# cars outside  $\leq$  capacity*

## Our goal

Modify refinement design of existing specification

→ improve understandability, maintainability, extensibility

# Example of Motivation

## Problem

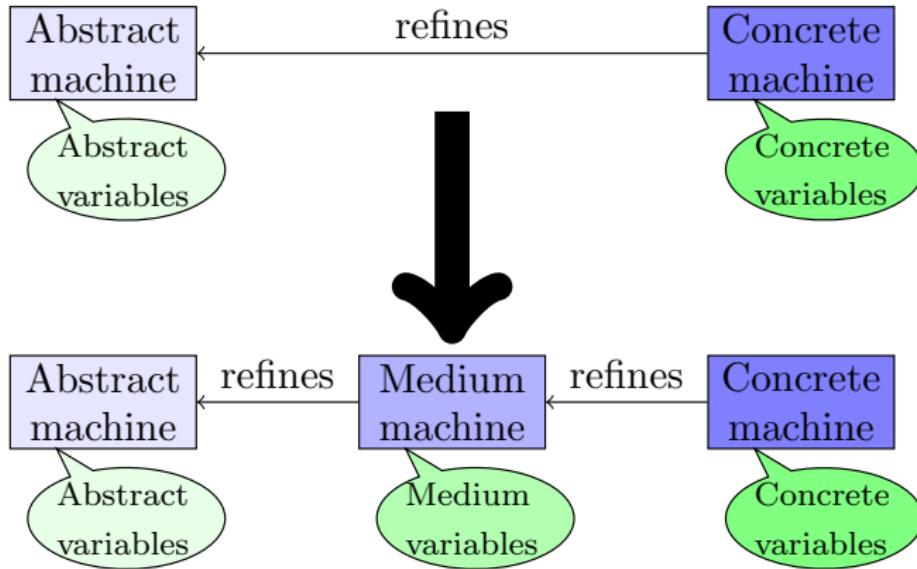
Sometimes we make refinements with many additional variables/invariants

- To specify several aspects in a step

## Solution

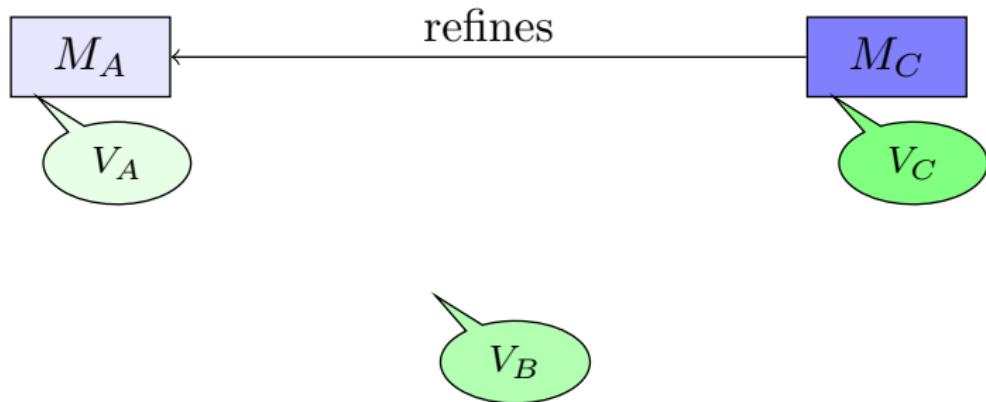
Refinement decomposition

# Refactoring of Refinement – Decomposition



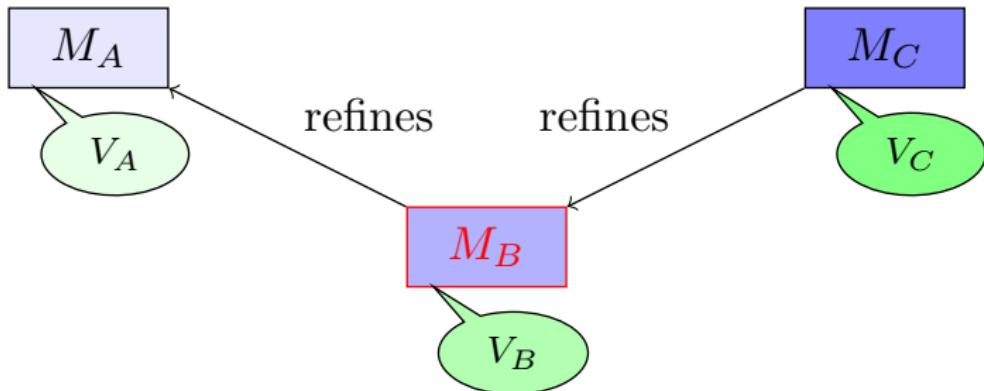
- Decomposition of refinement

# Goal of Refinement Decomposition



- Input
  - ▶ **Proved** machines  $M_A$  and  $M_C$
  - ▶ A set of variables  $V_B$  (subset of  $V_C$ , **slicing criteria**)

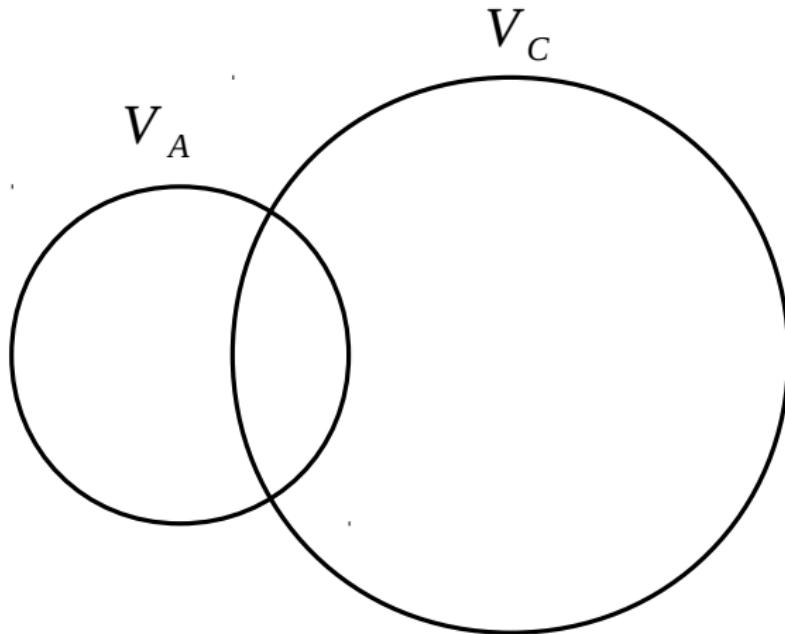
# Goal of Refinement Decomposition



- Input
  - ▶ **Proved** machines  $M_A$  and  $M_C$
  - ▶ A set of variables  $V_B$  (subset of  $V_C$ , **slicing criteria**)
- Output: Intermediate machine  $\textcolor{red}{M_B}$  such that
  - ▶  $M_C$  refines  $\textcolor{red}{M_B}$  and  $\textcolor{red}{M_B}$  refines  $M_A$
  - ▶  $\textcolor{red}{M_B}$  is specified in  $V_B$

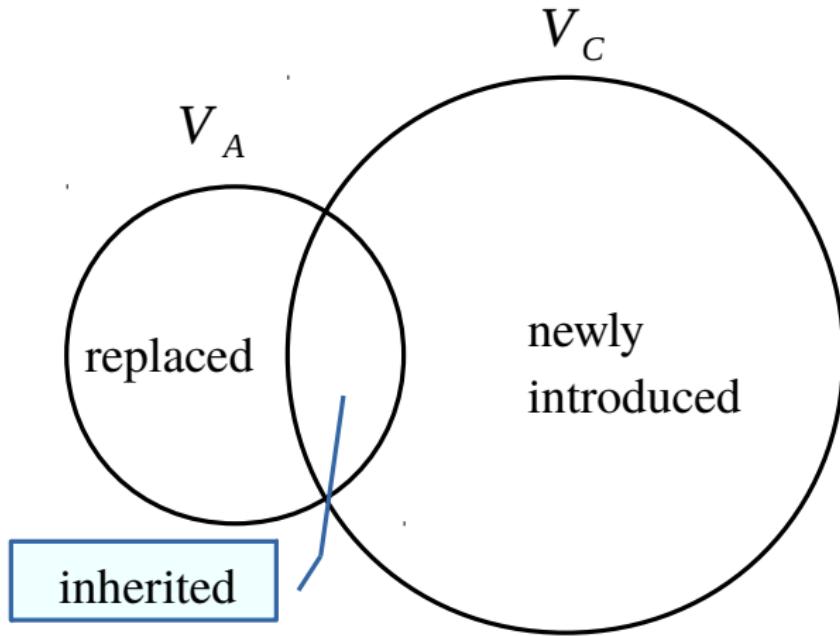
## Restriction on $V_B$

Variables in the machines:



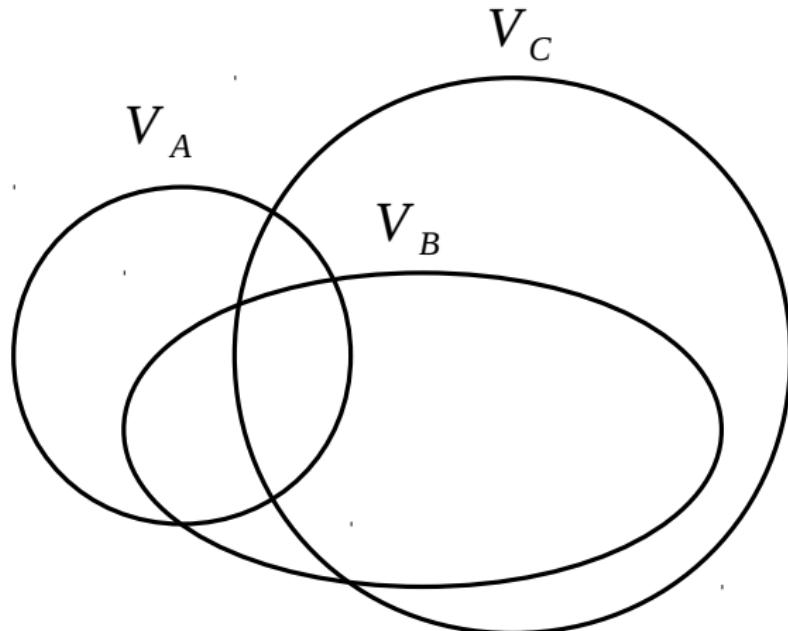
## Restriction on $V_B$

Variables in the machines:



## Restriction on $V_B$

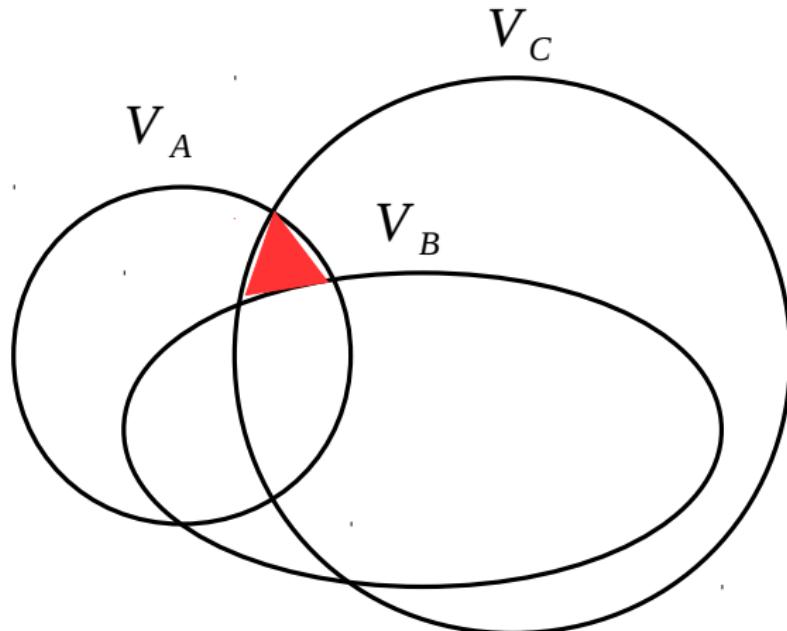
Variables in the machines:



Arbitrary  $V_B$ :

## Restriction on $V_B$

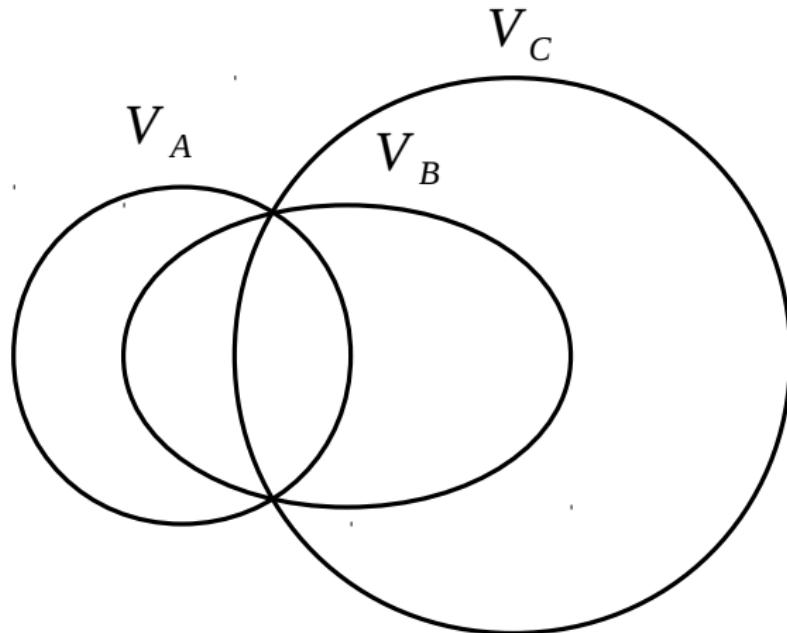
Variables in the machines:



Arbitrary  $V_B$ : Some variables are in  $V_A$  and  $V_C$  but not  $V_B$  !

## Restriction on $V_B$

Variables in the machines:



$\Rightarrow V_B$  should be a superset of  $V_A \cap V_C$

# Approach to Decomposing Refinement

1. Slicing from original machines  $M_C$  and  $M_A$

by finding specifications that can be expressed by  $V_B$

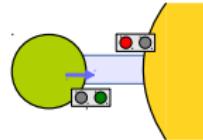
2. Mending for consistency

by providing complementary predicates  
to fill the gap originated from slicing

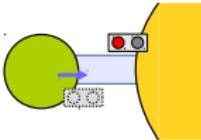
# Slicing

Find predicates that can be expressed by  $V_B$

e.g.,



$M_C (TL_{Island} \in V_C)$



$M_B (TL_{Island} \notin V_B)$

## Invariants

$TL_{Mainland} = \text{red} \vee TL_{Island} = \text{red}$

$TL_{Mainland} = \text{green} \Rightarrow n_{\rightarrow} = 0$

## Invariants

$TL_{Mainland} = \text{green} \Rightarrow n_{\rightarrow} = 0$

Event leave\_island

when

$TL_{Island} = \text{green}$

$1 \leq n_{Island}$

then

$n'_{Island} = n_{Island} - 1$

$n'_{\rightarrow} = n_{\rightarrow} + 1$

end

$\xrightarrow{\text{slice}}$

Event leave\_island

when

$1 \leq n_{Island}$

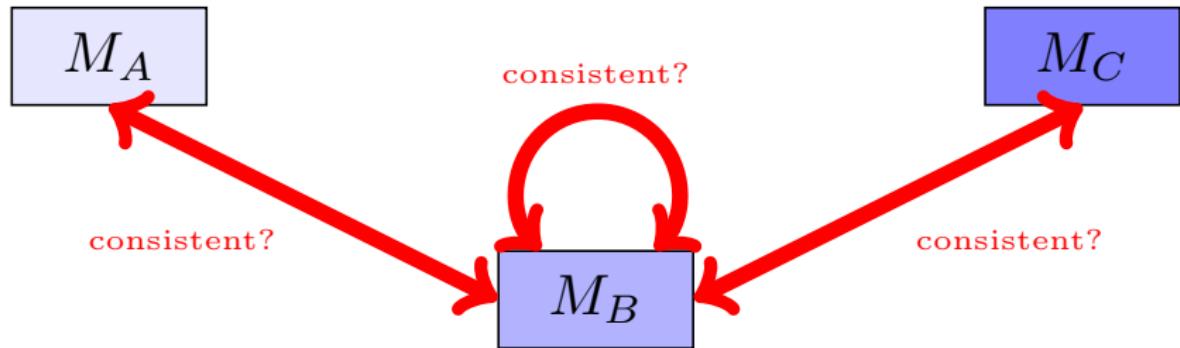
then

$n'_{Island} = n_{Island} - 1$

$n'_{\rightarrow} = n_{\rightarrow} + 1$

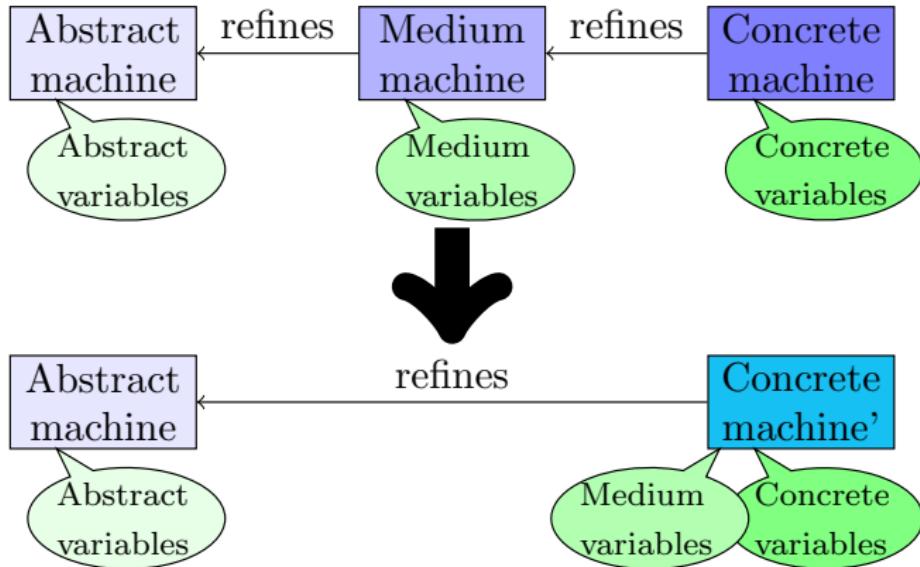
end

# Intermediate Machine is Not Always Consistent



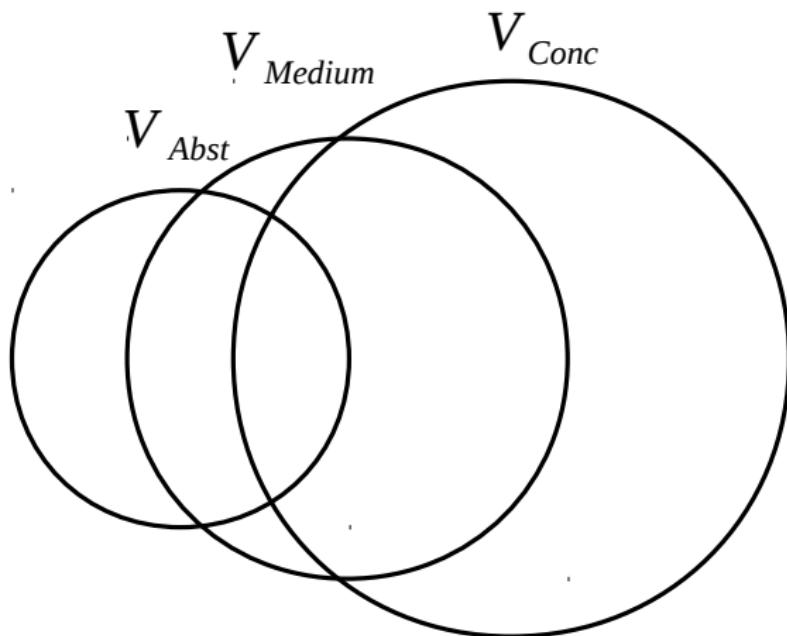
Sometimes we need to guarantee these **consistencies**

# Refactoring of Refinement – Composition

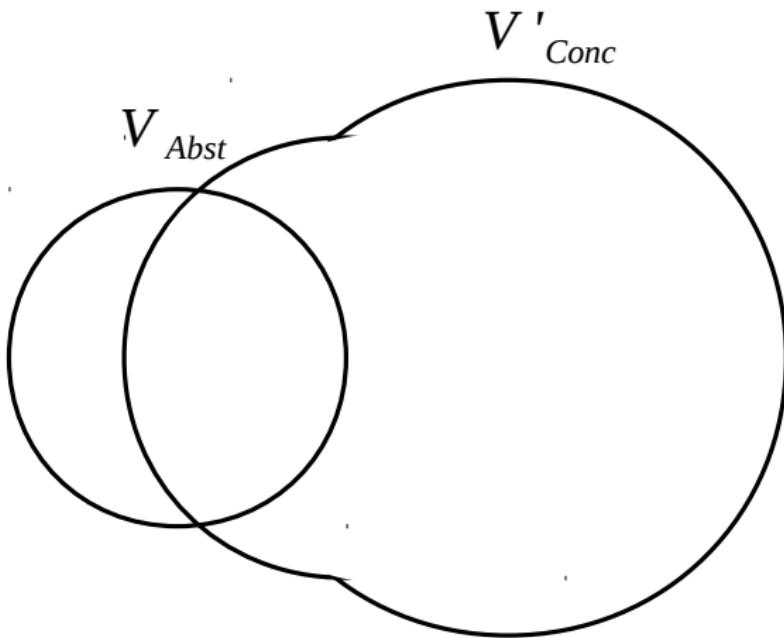


- Composition of refinement

# Refactoring of Refinement – Composition



# Refactoring of Refinement – Composition

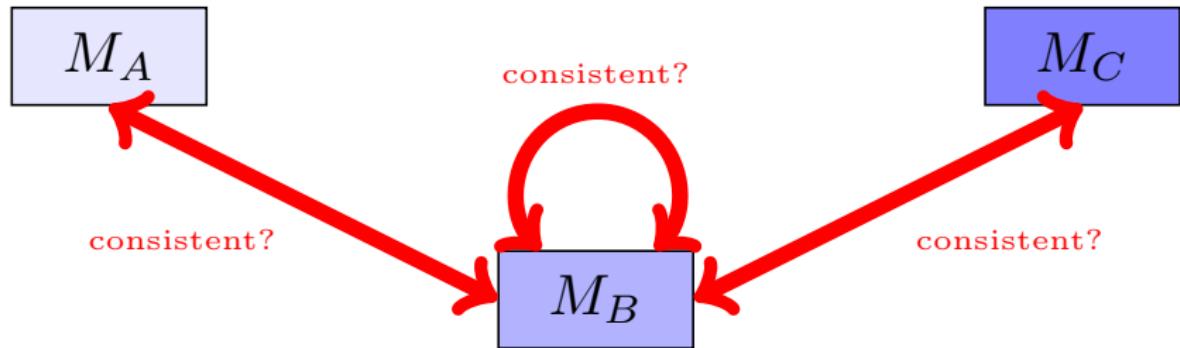


# SLICEANDMERGE: Implementation

The screenshot shows the Rodin IDE interface with the 'm1' model open. The top menu bar includes File, Edit, Navigate, Search, Project, Run, Window, and Help. The toolbar contains various icons for file operations like Open, Save, and Import. A 'Quick Access' button is in the top right. The main workspace displays a table with four columns: Element, Content, Special, and Comment. The table rows are categorized by element type (Invariants, Variables, Events, INITIALISATION, ML\_out, Parameters) indicated by icons in the first column.

| Element        | Content  | Special                | Comment   |
|----------------|--|------------------------|---|
| ■ Invariants   |  |                        |   |
| inv1           | $a \in N$  | not theorem            |   |
| inv2           | $b \in N$  | not theorem            |   |
| inv3           | $c \in N$  | not theorem            |   |
| inv4           | $n = a + b + c$  | not theorem            |   |
| inv5           | $a = 0 \vee c = 0$   | not theorem            | Gluing invariant connecting the concrete variables (a, b, c) to the abstract one (n)<br>The bridge is one-way |
| thm1           | $a + b + c \in N$  | theorem                |   |
| thm2           | $c > 0 \vee a > 0 \vee (a + b < d \wedge c = 0) \vee (0 < b \wedge a = 0)$ | theorem                |   |
| ■ Variables    |  |                        |   |
| a              |  |                        | number of cars on bridge going to island  |
| b              |  |                        | number of cars on island  |
| c              |  |                        | number of cars on bridge going to mainland  |
| ■ Events       |  |                        |   |
| INITIALISATION |  | not extended, ordinary |   |
| ML_out         |  | not extended, ordinary | leaving mainland  |
| Parameters     |  |                        |   |

# Intermediate Machine is Not Always Consistent



Sometimes we need to guarantee these **consistencies**

# Variables and Provability of Proof Obligations

Preservation of  $TL_{Mainland}$ 's property by "leaving from island"

|                           |  |
|---------------------------|--|
| Invariants                | $TL_{Mainland} = red \vee TL_{Island} = red$             |
| Guards                    | $TL_{Island} = green$                                    |
| Before-after predicates   | $\dots$  |
|                           | $n'_{\rightarrow} = n_{\rightarrow} + 1$                 |
| $\vdash$                  | $\vdash$   |
| Invariant after the event | $TL_{Mainland} = green \Rightarrow n'_{\rightarrow} = 0$ |

## Provability in {concrete, medium} machine

Concrete machine Provable.  $\because TL_{Mainland} \neq green$  from hypotheses

- Because of  $TL_{Island}$ -related predicates
  - either  $TL_{Mainland}$  or  $TL_{Island}$  is red
  - $TL_{Island}$  is green

Medium machine Not provable

- Because of lack of  $TL_{Island}$ -related predicates

# Mending by Adding Complementary Predicates

## Idea

- Original machines are consistent
  - $M_B$  lacks essential predicates for consistencies because of vocabulary limitations
- ↓
- Find essential predicates
  - Express them in vocabulary of  $V_B$  and mend  $M_B$  with them

## Ways of mending include:

- Heuristics such as extracting a predicate  $P$  from  $P \wedge Q$
- Analyzing the proof (as described in previous slide)
  - ① Trace the proof of original machine's consistency
  - ② Infer the complementary predicates from that
- Using *Craig's interpolation theorem*?

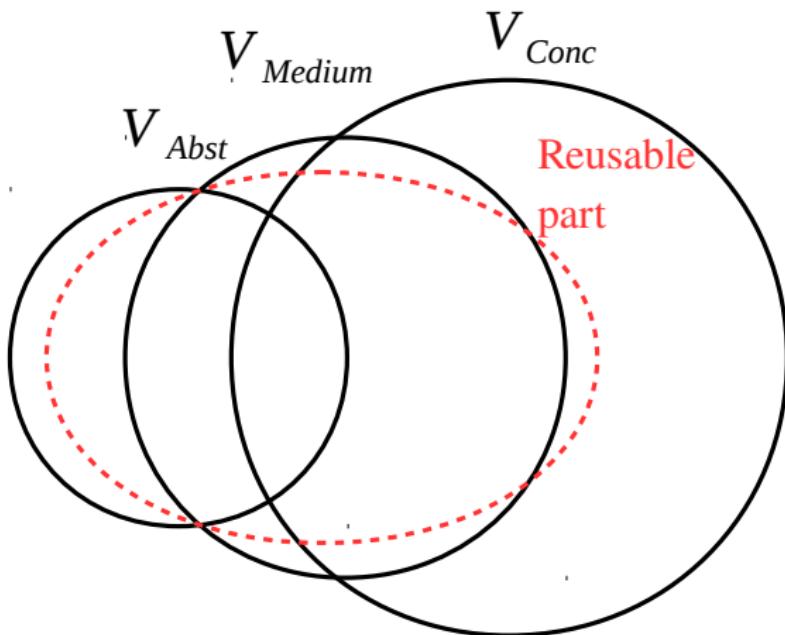


# Results

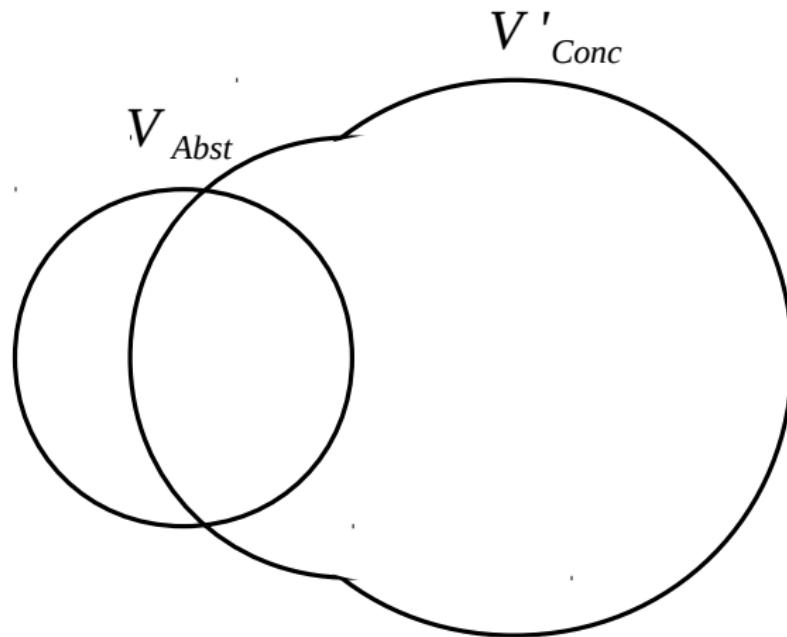
| original machines     | machine2 |     |     |     |  |
|-----------------------|----------|-----|-----|-----|--|
| intermediate machines | 1st      | 2nd | 3rd | 4th | Distributed introduction of variables/invariants |
| # variables           | +7       | +3  | +2  | +1  | +1   |
| # invariants          | 46       | 11  | 10  | 13  | 12   |
| # complement preds    | 12       | 13  | 10  | 0   |  |
| (# unique comp preds) | (5)      | (7) | (5) | (0) |  |

| original machines     | machine3 |      |      |     |                          |
|-----------------------|----------|------|------|-----|--------------------------|
| intermediate machines | 1st      | 2nd  | 3rd  | 4th | # found Complement preds |
| # variables           | -2+10    | -1+3 | -1+3 | +2  | +2                       |
| # invariants          | 72       | 7    | 17   | 14  | 34                       |
| # complement preds    | 20       | 20   | 15   | 0   |                          |
| (# unique comp preds) | (8)      | (9)  | (8)  | (0) |                          |

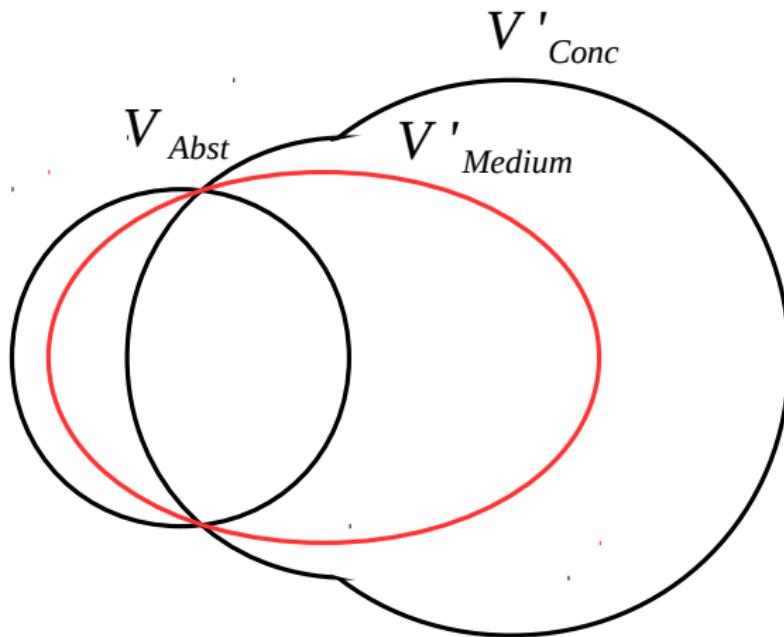
# Extracting Parts of Machines for Reuse



# Extracting Parts of Machines for Reuse



# Extracting Parts of Machines for Reuse



# Case Study: Extracting Authentication Parts of LAC

## Original model: location access controller

- Persons move between locations connected with turnstiles
- Persons are authorized to enter certain locations
- Persons insert their ID card in card readers on turnstiles
- Turnstiles communicate with a controller via messages
  - ▶ authentication, movement of a person, indicator light, ...

## New model: consoles, servers, and monitors

- Location have consoles with card readers and monitors.
- Authorized persons can login server by inserting ID card to the reader
- Controller try to find an unoccupied monitor in a room
- Consoles communicate with a controller via messages

# Concluding Remarks

## Summary

- Aiming at modify refinement structure of existing machines
- Slicing before mending / merging
- SLICEANDMERGE v1.0 coming soon?

## Future work

- Constructing systematic methods for mending
- *Planning refinement* for various use cases
- Finding other use cases
  - ▶ Planting specifications in other methods (e.g., VDM)?
  - ▶ ...

# Interpolation

## Theorem (Craig interpolation)

If a sequent **Hyp**  $\vdash$  **Goal** is provable,  
there exists an *interpolant* **I** such that

- All symbols of **I** occur in *both* **Hyp** and **Goal**
- **Hyp**  $\vdash$  **I** and **I**  $\vdash$  **Goal** are provable

## Example of interpolant

$$TL_{Mainland} = red \vee TL_{Island} = red$$

$$TL_{Island} = green$$

$$n'_{\rightarrow} = n_{\rightarrow} + 1$$

$\vdash$

$$TL_{Mainland} = green \Rightarrow n'_{\rightarrow} = 0$$

$\xleftarrow{\text{interpolates}}$

$$TL_{Mainland} \neq green$$

# Complementary Predicates and Interpolation

## Finding complementary predicates using interpolation

- ① Original consistency sequent  $\mathbf{Hyp} \vdash \mathbf{Goal}$  is given
- ② Apply **inference rules** to get another sequent  $\mathbf{Hyp}' \vdash \mathbf{Goal}'$  such that  $\mathbf{Goal}'$  can be expressed using  $V_B$
- ③ Obtain an interpolant  $I'$  of  $\mathbf{Hyp}' \vdash \mathbf{Goal}'$ 
  - ▶  $\mathbf{Hyp}' \vdash I', I' \vdash \mathbf{Goal}'$  //  $I'$  is an essence
  - ▶  $I'$  can be expressed in  $V_B$
- ④ Find complementary predicates from  $I'$

# Complementary Predicates and Interpolation

## Finding complementary predicates using interpolation

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$$TL_{Mainland} = red \vee TL_{Island} = red$$

$$TL_{Island} = green$$

$$n'_{\rightarrow} = n_{\rightarrow} + 1$$

⊤

$$TL_{Mainland} = green \Rightarrow n'_{\rightarrow} = 0$$

# Complementary Predicates and Interpolation

## Finding complementary predicates using interpolation

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- ④ Find complementary predicates from  $I'$

$$TL_{Mainland} = red \vee TL_{Island} = red$$

$$TL_{Island} = green$$

$\vdash$

$$TL_{Mainland} = green \Rightarrow n'_{\rightarrow} = 0$$

$$\neg n'_{\rightarrow} = n_{\rightarrow} + 1$$

# Complementary Predicates and Interpolation

## Finding complementary predicates using interpolation

- ① Original consistency sequent  $\mathbf{Hyp} \vdash \mathbf{Goal}$  is given
- ② Apply **inference rules** to get another sequent  $\mathbf{Hyp}' \vdash \mathbf{Goal}'$  such that  $\mathbf{Goal}'$  can be expressed using  $V_B$
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  - ▶  $I'$  can be expressed in  $V_B$
- ④ Find complementary predicates from  $I'$

$$TL_{Mainland} = red \vee TL_{Island} = red$$

$$TL_{Island} = green$$

$$\vdash \frac{TL_{Mainland} = green \Rightarrow n'_\rightarrow = 0 \quad \neg n'_\rightarrow = n_\rightarrow + 1}{TL_{Mainland} \neq green} \text{ interpolates}$$